

## [§4.8] Algorithms

### 3.1 Introduction & Notation

- An **algorithm** is a list of step-by-step instructions on **how to** accomplish a task. Required properties include:
  - Input/Output -- An algorithm takes some input and produces an output.
  - Precision -- Each step is clearly indicated and unambiguous.
  - Finiteness -- The algorithm stops after a finite number of steps.

*Example:* Find the maximum of three numbers, a, b, and c

- Non-algorithm  
"Compare the three (input) numbers and return (or output) the largest one."
- Algorithm
  - Step 1: Let x (another number) to have the value of a.
  - Step 2: If  $b > x$ , then set x's value to b.
  - Step 3: If  $c > x$ , then set x's value to c.
  - Step 4: Return x.
- Algorithms are often written as **pseudocode**, which is sort of the middle ground between a formal programming language and plain English.
- Some of the common elements of pseudocode:
  - Each algorithm starts with "**procedure** name(input values)".
  - There must be at least one "**return** value" statement.  
NOTE: A procedure can only return ONE value only.
  - *Operators* are:
    - $+$ ,  $-$ ,  $*$ ,  $/$ , mod
    - $<$ ,  $\leq$ ,  $>$ ,  $\geq$
    - $=$ ,  $\neq$
  - The *assignment operator* is  $:=$  (e.g. " $x := 2$ ").
- The basic constructs are
  - "**if** cond **then** action1 **else** action2"
  - "**for** var  $:=$  initial\_value **to** final value **do** action"
  - "**while** cond **do** action"
  - "**repeat** action **until** cond"
  - "**begin** statement<sub>1</sub> .. statement<sub>n</sub> **end**"
- *Example:* Algorithm *max* which finds the maximum of three numbers, a, b, and c.

Input: Three numbers a, b, c

Output: A number which is the largest of the tree input numbers

```
procedure max(a, b, c)
1. x := a           // x is a local/temporary variable
2. if b > x then
3.   x := b         // update x
4. if c > x then
5.   x := c         // update x
6. return x
```

## 3.2 Example Algorithms

1. Algorithm *find\_largest* which finds the largest number in the sequence  $s_1, \dots, s_n$ .

Input: A list of integers  $s = \{s_1, \dots, s_n\}$ , and the length  $n$  of the list

Output: A number which is the largest in  $s$ .

2. Algorithm *Div* which calculate, for a given integer ( $n$ ) and a positive integer ( $d$ ), the quotient ( $q$ ) and remainder ( $r$ ) where  $q$  and  $r$  are integers. Division is done by repeated subtraction, not by arithmetic division.

Input:  $n$  (an integer), and  $d$  (a positive integer)

Output:  $q$  and  $r$  (integers)

**procedure** Div( $n$ ,  $d$ ,  $q$ ,  $r$ )

3. Algorithm *is\_even* which tests whether a positive integer  $m$  is even

Input: A positive integer  $m$

Output: **true** if  $m$  is even; **false** if  $m$  is odd

## 3.3 The Euclidean Algorithm

- An algorithm which finds the **greatest common divisor (gcd)** of two integers.

*Example:*  $\text{gcd}(30, 105) = 15$

divisor of 30	
divisor of 105	
common	

- **"divides"**
  - If  $a$ ,  $b$  and  $q$  are integers (where  $b \neq 0$ ) satisfying  $a = bq$ , we say that "**b divides a**", noted  $b \mid a$ . The value  $q$  is called the **quotient**, and  $b$  a **divisor** of  $a$ .

*Examples:*  $15 \mid 30$ , and the quotient is 2.

- Properties of divisors:

Let  $m$ ,  $n$  and  $c$  be integers.

1. If  $c \mid m$  and  $c \mid n$ , then  $c \mid (m + n)$
2. If  $c \mid m$  and  $c \mid n$ , then  $c \mid (m - n)$
3. If  $c \mid m$ , then  $c \mid mn$

- **Theorem:** For all integers  $a, b, q$  and  $r$ , such that  $a = bq + r$ , and  $a \geq 0, b > 0$  and  $0 \leq r < b$ , then  $\gcd(a, b) = \gcd(b, r)$ .

First, let's check some cases:

- $105 = 30 * 3 + 15 \rightarrow \gcd(105, 30) = \gcd(30, 15) = 15$
- $504 = 396 * 1 + 108 \rightarrow \gcd(504, 396) = \gcd(396, 108) = 36$

Proof:

Let  $C1$  be the set of common divisors of  $a$  and  $b$ , and  $C2$  be the set of common divisors of  $b$  and  $r$ . We show  $C1 \supseteq C2$ , and  $C2 \supseteq C1$ . Then we can conclude  $C1 = C2$ .

1) Show  $C1 \supseteq C2$

Let  $c \in C1$  be a common divisor of  $a$  and  $b$ , that is,  $c \mid a$  and  $c \mid b$ .

From the third property of divisors, we get that  $c \mid bq$ .

Then from the second property of divisors, we get that  $c \mid (a - bq) = r$ . ... (A)

From (A) and hypothesis, we have that  $c \mid b$  and  $c \mid r$ . Therefore,  $c$  is a common divisor of  $b$  and  $r$ , that is,  $c \in C2$ .

2) Show  $C2 \supseteq C1$

Let  $c \in C2$  be a common divisor of  $b$  and  $r$ , that is,  $c \mid b$  and  $c \mid r$ .

From the third property of divisors, we get that  $c \mid bq$ .

Then from the first property of divisors, we get that  $c \mid (bq + r) = a$ . ... (B)

From (B) and hypothesis,  $c \mid a$  and  $c \mid b$ . Therefore,  $c$  is a common divisor of  $a$  and  $b$ , that is,  $c \in C1$ .

From 1) and 2), we showed that all common divisor of  $a$  and  $b$  are also common divisors of  $b$  and  $r$ . In particular, this is true for the greatest common divisor, i.e.,  $\gcd(a, b) = \gcd(b, r)$ .

Algorithm

```

procedure gcd(a, b)
1. if a < b, then
2.   swap(a, b)      // make a the larger of the two
3. while b != 0 do
4.   begin
5.     r := a mod b
6.     a := b
7.     b := r
8.   end
9. return a

```

iteration		a	b	r
1	$\gcd(105, 30)$	105	30	15
2				
3				

- *Example:*  $\gcd(110, 273)$  -- Section 3.3, question #2

iteration		a	b	r
1				
2				
3				