

Lecture Note #8 (Recurrence Relations)

Exercises

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(1) Iteration Method

More Examples:

a. $T(n) = -3 \cdot T(n-1)$, $T(0) = 2$

Solution: By iteration.

$$\begin{aligned} T(n) &= (-3) \cdot T(n-1) \\ &= (-3) \cdot [(-3) \cdot T(n-2)] \\ &= (-3)^2 \cdot T(n-2) \\ &= (-3)^2 \cdot [(-3) \cdot T(n-3)] \\ &= (-3)^3 \cdot T(n-3) \\ &= \dots \\ &= (-3)^k \cdot T(n-k) \\ &= \dots \\ &= (-3)^n \cdot T(0) \dots \text{because } n-k=0, \text{ which gives } k=n \\ &= (-3)^n \cdot (2) \end{aligned}$$

So the closed form is $C(n) = (-3)^n \cdot (2)$ for all integer $n \geq 0$.

b. $T(n) = T(n-1) + 1$, $T(0) = 0$

Solution: By iteration.

$$\begin{aligned} T(n) &= T(n-1) + 1 \\ &= [T(n-1) + 1] + 1 \\ &= T(n-1) + 2 \\ &= \dots \\ &= T(n-k) + k \\ &= \dots \\ &= T(0) + n \\ &= n \end{aligned}$$

So the closed form is $C(n) = n$ for all integer $n \geq 0$.

c. $T(n) = 2 \cdot T(n-1) + 3$, $T(0) = 0$

Solution: By iteration.

$$\begin{aligned} T(n) &= 2 \cdot T(n-1) + 3 \\ &= 2 \cdot [2 \cdot T(n-2) + 3] + 3 \\ &= 2^2 \cdot T(n-2) + (2) \cdot 3 + 3 \\ &= 2^2 \cdot [2 \cdot T(n-3) + 3] + (2) \cdot 3 + 3 \end{aligned}$$

$$\begin{aligned}
&= 2^3 \cdot T(n-3) + 2^2 \cdot 3 + 2^1 \cdot 3 + 2^0 \cdot 3 \\
&= \dots \\
&= 2^k \cdot T(n-k) + \sum_{i=0}^{k-1} 3 \cdot 2^i \\
&= 2^k \cdot T(n-k) + 3 \cdot \sum_{i=0}^{k-1} 2^i \\
&= \dots \\
&= 2^n \cdot T(0) + 3 \cdot \sum_{i=0}^{n-1} 2^i \\
&= 2^n \cdot T(0) + 3 \cdot \frac{2^n - 1}{2 - 1} \\
&= 2^n \cdot 0 + 3 \cdot (2^n - 1) \\
&= 3 \cdot (2^n - 1)
\end{aligned}$$

So the closed form is $C(n) = 3 \cdot (2^n - 1)$ for all integer $n \geq 0$.

d. $T(n) = T(n/2) + 1, T(1) = 0$

Solution: By iteration.

$$\begin{aligned}
T(n) &= T\left(\frac{n}{2}\right) + 1 \\
&= \left[T\left(\frac{n}{2^2}\right) + 1\right] + 1 \\
&= T\left(\frac{n}{2^2}\right) + 2 \\
&= \left[T\left(\frac{n}{2^3}\right) + 1\right] + 2 \\
&= T\left(\frac{n}{2^3}\right) + 3 \\
&= \dots \\
&= T\left(\frac{n}{2^k}\right) + k \\
&= \dots \\
&= T(1) + \log_2 n \dots \text{because } 2^k = n, \text{ which gives } k = \log_2 n \\
&= \log_2 n
\end{aligned}$$

So the closed form is $C(n) = \log_2 n$ for all integers $n \geq 1$.

Proving the correctness of the derived formula

* **Example 2**: Show if $T(1), T(2), \dots, T(n)$ is the sequence defined by

$$T(n) = -3 \cdot T(n-1) \text{ for all integers } n \geq 1, \text{ and } T(0) = 2$$

Then the closed form is $T(n) = 2 \cdot (-3)^n$, for all $n \geq 0$

Proof: By induction.

Basis step ($n = 0$):

- Recurrence: $T(0) = 2$, as given.
- Closed form: $T(0) = 2 \cdot (-3)^0 = 2 \cdot 1 = 2$.

Therefore, the closed form holds for $n = 0$ (A)

Inductive step:

Assume the closed form holds for $n-1$, that is, $T(n-1) = 2 \cdot (-3)^{n-1}$, for all integers $n \geq 1$.
Show the closed form holds for n as well, that is, $T(n) = 2 \cdot (-3)^n$

$$\begin{aligned} T(n) &= (-3) \cdot T(n-1) \dots \text{by recurrence relation} \\ &= (-3) \cdot [2 \cdot (-3)^{n-1}] \dots \text{by inductive hypothesis} \\ &= 2 \cdot (-3)^n \dots \text{as to be shown ... (B)} \end{aligned}$$

By (A) and (B), we can conclude that the closed form holds for all integers $n \geq 0$. QED.